

## SECTION - A

## GENERAL APTITUDE

Q.1 Given below are two statements I and II and two conclusions I and II:

**Statement:**

- I. All bacteria are microorganisms.
- II. All pathogens are microorganisms.

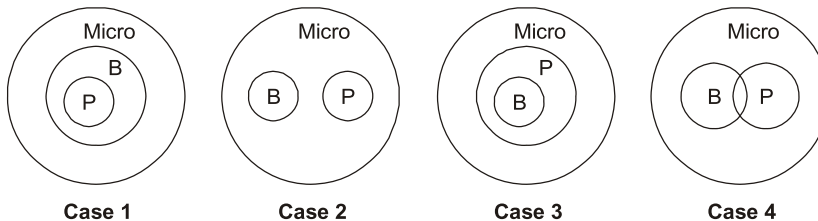
**Conclusions:**

- I. Some pathogens are bacteria.
- II. All pathogens are not bacteria.

Based on the above statements and conclusions, which one of the following options is logically CORRECT?

- (a) Only conclusion I is correct
- (b) Either conclusion I or II is correct.
- (c) Only conclusion II is correct
- (d) Neither conclusion I nor II is correct.

Ans. (d)



None of the two conclusions will satisfy all the 4 cases.

End of Solution

Q.2 There are five bags each containing identical sets of ten distinct chocolates. One chocolate is picked from each bag.

The probability that at least two chocolates are identical is \_\_\_\_\_.

- (a) 0.6976
- (b) 0.3024
- (c) 0.4235
- (d) 0.8125

Ans. (a)

10 distinct chocolates, 5 bags

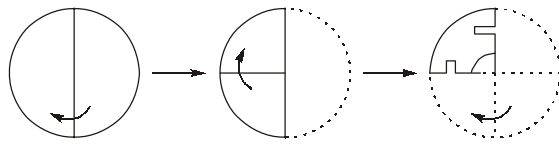
$$P(\text{atleast 2}) = 1 - P(\text{exactly 1})$$

$$P(\text{exactly 1}) = P(\text{all different}) = \frac{10 \times 9 \times 8 \times 7 \times 6}{10^5}$$

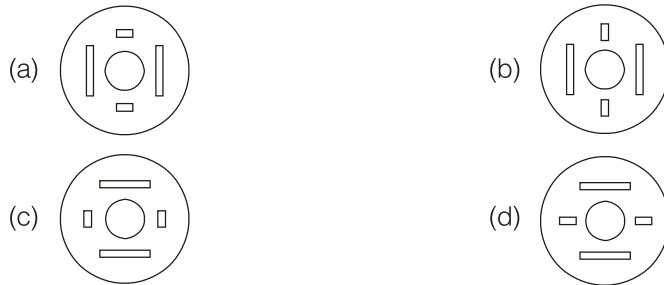
$$\begin{aligned} P(\text{atleast 2}) &= 1 - \frac{10 \times 9 \times 8 \times 7 \times 6}{10^5} \\ &= 1 - 0.3024 = 0.6976 \end{aligned}$$

End of Solution

Q.3



A circular sheet of paper is folded along the lines in the directions shown. The paper, after being punched in the final folded state as shown and unfolded in the reverse order of folding, will look like \_\_\_\_\_.



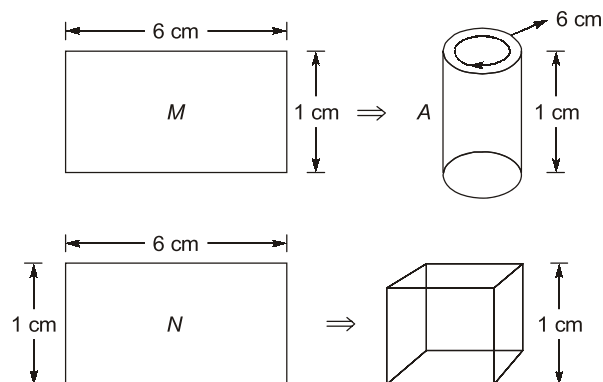
Ans. (c)

End of Solution

Q.4 We have 2 rectangular sheets of paper, M and N, of dimensions 6 cm × 1 cm each. Sheet M is rolled to form an open cylinder by bringing the short edges of the sheet together. Sheet N is cut into equal square patches and assembled to form the largest possible closed cube. Assuming the ends of the cylinder are closed, the ratio of the volume of the cylinder to that of the cube is \_\_\_\_\_.

- (a)  $\frac{9}{\pi}$  (b)  $3\pi$   
 (c)  $\frac{\pi}{2}$  (d)  $\frac{3}{\pi}$

Ans. (a)



$$\text{Volume of cylinder} = \pi r^2 h$$

Now,  $2\pi r = 6$  (figure A)

$$r = \frac{3}{\pi}$$

$$\text{Volume of cylinder} = \pi \times \frac{3}{\pi} \times \frac{3}{\pi} \times 1 = \frac{9}{\pi}$$

$$\text{Volume of cube} = (1)^3$$

$$\text{Ratio} = \frac{\frac{9}{\pi}}{1} = \frac{9}{\pi}$$

End of Solution

**Q.5** Consider the following sentences:

- (i) Everybody in the class is prepared for the exam.
- (ii) Babu invited Danish to his home because he enjoys playing chess,

Which of the following is the CORRECT observation about the above two sentences?

- (a) (i) is grammatically incorrect and (ii) is unambiguous
- (b) (i) is grammatically correct and (ii) is ambiguous
- (c) (i) is grammatically incorrect and (ii) is ambiguous
- (d) (i) is grammatically correct and (ii) is unambiguous

**Ans. (b)**

Everybody is singular and takes singular verb is, first statement is correct and the purpose of invitation is unclear so second statement is vague and that sounds ambiguous.

End of Solution

**Q.6** Some people suggest anti-obesity measures (AOM) such as displaying calorie information in restaurant menus. Such measures sidestep addressing the core problems that cause obesity: poverty and income inequality.

Which one of the following statements summarizes the passage?

- (a) The proposed AOM addresses the care problems that cause obesity.
- (b) If obesity reduces, poverty will naturally reduce, since obesity causes poverty.
- (c) AOM are addressing the problem superficially.
- (d) AOM are addressing the core problems and are likely to succeed.

**Ans. (c)**

Superficially is the deciding key word which means apparently/seemingly.

End of Solution

**Q.7** \_\_\_\_\_ is to *surgery* as *writer* is to \_\_\_\_\_

Which one of the following options maintains a similar logical relation in the above sentence?

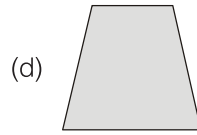
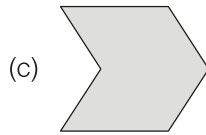
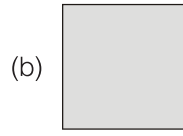
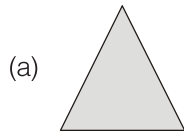
- (a) Plan, outline
- (b) Hospital, library
- (c) Doctor, book
- (d) Medicine, grammar

**Ans. (c)**

Doctor performs surgery as writer writes book.

End of Solution

**Q.8** A polygon is convex if, for every pair of points, P and Q belonging to the polygon, the line segment PQ lies completely inside or on the polygon.  
Which one of the following is NOT a convex polygon?



**Ans. (c)**

**End of Solution**

**Q.9**

Items	Cost (₹)	Profit%	Marked price (₹)
P	5400	–	5860
Q	–	25	10000

Details of prices of two items P and Q are presented in the above table. The ratio of cost of item P to cost of item Q is 3 : 4. Discount is calculated as the difference between the marked price and the selling price. The profit percentage is calculated as the ratio of the difference between selling price and cost, to the cost

$$\left( \text{Profit\%} = \frac{\text{Selling price} - \text{Cost}}{\text{Cost}} \times 100 \right).$$

The discount on item Q, as a percentage of its marked price, is \_\_\_\_\_.

- (a) 25 (b) 10  
(c) 5 (d) 12.5

**Ans. (b)**

$$\Rightarrow CP_P = 5400, \quad MP_P = 5860$$

$$\frac{5400 \times 4}{5} = CP_Q$$

$$CP_Q = 7200$$

$$SP_Q = ?$$

$$\text{Profit \%} = \frac{\text{Profit}}{\text{CP}} \times 100 \text{ [for Q]}$$

$$25 = \left( \frac{SP - 7200}{7200} \right) 100$$

$$SP_Q = 9000$$

$$\text{Discount } Q = MP_Q - SP_Q = 1000$$

$$\text{Discount \%} = \frac{1000}{10000} \times 100 = 10\%$$

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**End of Solution**

**Q.10** The ratio of boys to girls in a class is 7 to 3.  
Among the options below, an acceptable value for the total number of students in the class is:

- (a) 37 (b) 50  
(c) 21 (d) 73

**Ans. (b)**

$$\text{Boys} = 7x, \quad \text{Girls} = 3x$$

$$\text{Total number of students} = 10x$$

Now,  $7x$  and  $3x$  should be integers.

Hence,  $x$  should be integer.

So  $10x$  should be an integer.

Only 50 satisfies the above constraints.

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**End of Solution**



**SECTION - B****TECHNICAL**

**Q.1** Consider the following C code segment:

```
a = b + c;  
e = a + 1;  
d = b + c;  
f = d + 1;  
g = e + f;
```

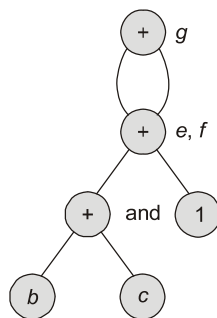
In a compiler, this code segment is represented internally as a directed acyclic graph (DAG). The number of nodes in the DAG is \_\_\_\_\_.

**Ans. (6)**

```
a = b + c;  
e = a + 1;  
d = b + c;  
f = d + 1;  
g = e + f;
```

Using common sub-expression elimination this code will become.

```
a = b + c  
e = a + 1  
d = a  
f = e  
g = e + e
```



Number of nodes = 6

**End of Solution**

**Q.2** Consider the two statements:

$S_1$  : There exist random variables X and Y such that  
 $(E[X - E(X)](Y - E(Y)))^2 > \text{Var}[X] \text{Var}[Y]$

$S_2$  : For all random variables X and Y  
 $\text{Cov}[X, Y] = E[(X - E[X]) (Y - E[Y])]$

Which one of the following choices is correct?

- (a) Both  $S_1$  and  $S_2$  are false.                      (b) Both  $S_1$  and  $S_2$  are true.  
(c)  $S_1$  is true, but  $S_2$  is false.                      (d)  $S_1$  is false, but  $S_2$  is true.

**Ans. (a)**

$S_2$  :  $\text{Cov}(x, y) = E\{(x - \bar{x})(y - \bar{y})\}$  is false

**Case-I:** If  $x > \bar{x}$  and  $y > \bar{y}$  then above is true.

**Case-II:** If  $x < \bar{x}$  and  $y < \bar{y}$  then above is true.

**Case-III:** If  $x > \bar{x}$  but  $y < \bar{y}$  then above is false.

**Case-IV:** If  $x < \bar{x}$  but  $y > \bar{y}$  then above is false.

$\therefore$  Given expression is not always true. So we can conclude that it is false.

$S_1$  : It is obviously false.

$\therefore$  True statement is  $[E\{(x - \bar{x})(y - \bar{y})\}]^2 < \text{var}(x) \cdot \text{Var}(y)$

So both  $S_1$  and  $S_2$  are false.

**End of Solution**

**Q.3** Consider the following two statements:

$S_1$  : Destination MAC address of an ARP reply is a broadcast address.

$S_2$  : Destination MAC address of an ARP request is a broadcast address.

Which one of the following choices is correct?

- (a)  $S_1$  is false and  $S_2$  is true.                      (b) Both  $S_1$  and  $S_2$  are false.  
(c) Both  $S_1$  and  $S_2$  are true,                      (d)  $S_1$  is true and  $S_2$  is false.

**Ans. (a)**

ARP request is broadcasting.

ARP reply is unicasting.

**End of Solution**

**Q.4** Three processes arrive at time zero with CPU bursts of 16, 20 and 10 milliseconds. If the scheduler has prior knowledge about the length of the CPU bursts, the minimum achievable average waiting time for these three processes in a non-preemptive scheduler (round to nearest integer) is \_\_\_\_\_ milliseconds.

**Ans. (12)**

WT	CT	P. No.	AT	BT
0	10	A	0	16
26	46	B	0	20
10	26	C	0	10

C	A	B
0	10	26

$$\begin{aligned} \text{Average waiting time} &= \frac{0+26+10}{3} \\ &= \frac{36}{3} = 12 \end{aligned}$$

**End of Solution**

**Q.5** Consider the following matrix:

$$\begin{pmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

The largest eigenvalue of the above matrix is \_\_\_\_\_.

**Ans. (3)**

$$A = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

$$|A - \lambda I| = 0 \Rightarrow \begin{vmatrix} -\lambda & 1 & 1 & 1 \\ 1 & -\lambda & 1 & 1 \\ 1 & 1 & -\lambda & 1 \\ 1 & 1 & 1 & -\lambda \end{vmatrix} = 0$$

$$\begin{vmatrix} -\lambda+3 & 1 & 1 & 1 \\ -\lambda+3 & -\lambda & 1 & 1 \\ -\lambda+3 & 1 & -\lambda & 1 \\ -\lambda+3 & 1 & 1 & -\lambda \end{vmatrix} = 0$$

$$C_1 \leftarrow C_1 + C_2 + C_3 + C_4$$



$$\Rightarrow (-\lambda + 3) \begin{vmatrix} 1 & 1 & 1 & 1 \\ 1 & -\lambda & 1 & 1 \\ 1 & 1 & -\lambda & 1 \\ 1 & 1 & 1 & -\lambda \end{vmatrix} = 0$$

$$R_2 \leftarrow R_2 - R_1, R_3 \leftarrow R_3 - R_1, R_4 \leftarrow R_4 - R_1$$

$$\Rightarrow (-\lambda + 3) \begin{vmatrix} 1 & 1 & 1 & 1 \\ 0 & -\lambda - 1 & 0 & 0 \\ 0 & 0 & -\lambda - 1 & 0 \\ 0 & 0 & 0 & -\lambda - 1 \end{vmatrix} = 0$$

$$\text{or } (-\lambda + 3) (-\lambda - 1)^3 = 0$$

$$\Rightarrow \lambda = 3, -1, -1, -1,$$

So maximum eigen value is  $\lambda = 3$ .

**End of Solution**

**Q.6** Consider the following grammar (that admits a series of declarations, followed by expressions) and the associated syntax directed translation (SDT) actions, given as pseudo-code:

$P \rightarrow D^* E^*$

$D \rightarrow \text{int ID}$  {record that ID.lexeme is of type int}

$D \rightarrow \text{bool ID}$  {record that ID.lexeme is of type bool}

$E \rightarrow E_1 + E_2$  {check that  $E_1.type = E_2.type = \text{int}$ ; set  $E.type := \text{int}$ }

$E \rightarrow !E_1$  {check that  $E_1.type = \text{bool}$ ; set  $E.type := \text{bool}$ }

$E \rightarrow \text{ID}$  {set  $E.type := \text{int}$ }

With respect to the above grammar, which one of the following choices is correct?

- (a) The actions can be used to type-check syntactically correct integer variable declarations and integer expressions.
- (a) The actions can be used to type-check syntactically correct boolean variable declarations and boolean expressions.
- (c) The actions will lead to an infinite loop.
- (d) The actions can be used to correctly type-check any syntactically correct program.

**Ans. (a)**

1.  $P \rightarrow D^* E^*$

2.  $D \rightarrow \text{int ID}$  {record that ID.lexeme is of type int}

3.  $D \rightarrow \text{bool ID}$  {record that ID.lexeme is of type bool}

4.  $E \rightarrow E_1 + E_2$  {check that  $E_1.type = E_2.type = \text{int}$ ; set  $E.type = \text{int}$ }

5.  $E \rightarrow !E_1$  {check that  $E_1.type = \text{bool}$ ; set  $E.type = \text{bool}$ }

6.  $E \rightarrow \text{ID}$  {set  $E.type = \text{int}$ }

Rules 2 and 3 are used for entry into the symbol table. Rule 4 is used for type checking of the integer expression. But, in rule 6 only int type is set.

Hence answer is option (a).

**End of Solution**

**Q.7** Define  $R_n$  to be the maximum amount earned by cutting a rod of length  $n$  meters into one or more pieces of integer length and selling them. For  $i > 0$ , let  $p[i]$  denote the selling price of a rod whose length is  $i$  meters. Consider the array of prices:

$$p[1] = 1, p[2] = 5, p[3] = 8, p[4] = 9, p[5] = 10, p[6] = 17, p[7] = 18$$

Which of the following statements is/are correct about  $R_7$ ?

- (a)  $R_7 = 19$
- (b)  $R_7 = 18$
- (c)  $R_7$  is achieved by three different solutions.
- (d)  $R_7$  cannot be achieved by a solution consisting of three pieces.

**Ans. (b, c)**

$R_7$  : Maximum amount earned by cutting rod of length '7' into 1, 2, 3, 4, 5, 6, 7 pieces (whichever way is maximum)

$$p[7] = 18$$

$$p[6] + p[1] = 17 + 1 = 18$$

(Also,  $R_7$  is achieved by 3 diff. solution)

$$p[5] + p[2] = 10 + 5 = 15$$

$$p[4] + p[3] = 9 + 8 = 17$$

$$p[5] + p[1] + p[1] = 10 + 1 + 1 = 12$$

$$p[4] + p[2] + p[1] = 9 + 5 + 1 = 15$$

$$p[4] + p[1] + p[1] + p[1] = 9 + 3(1) = 12$$

$$p[3] + p[4] = 17$$

$$p[3] + p[3] + p[1] = 8 + 8 + 1 = 17$$

$$p[3] + p[2] + p[2] = 8 + 5 + 5 = 18$$

( $R_7$  is achieved by 3 diff. pieces)

$$p[3] + 4 p[1] = 8 + 4 = 12$$

$$p[2] + p[5] = 15$$

$$p[2] + p[4] + p[1] = 5 + 9 + 1 = 15$$

$$p[2] + p[3] + p[2] = 18$$

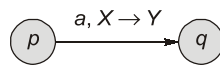
$$p[2] + 5 p[1] = 5 + 5 = 10$$

$$p[2] + p[2] + 3 p[1] = 5 + 5 + 3 = 13$$

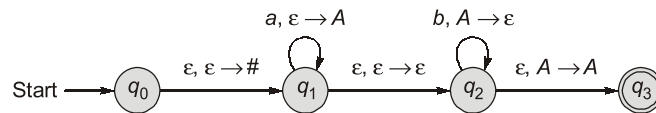
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**End of Solution**

**Q.8** In a pushdown automaton  $P = (Q, \Sigma, \Gamma, \delta, q_0, F)$ , a transition of the form,



where  $p, q, \in Q$ ,  $a \in \Sigma \cup \{\epsilon\}$ , and  $X, Y \in \Gamma \cup \{\epsilon\}$ , represents  $(q, Y) \in \delta(p, a, X)$   
 Consider the following pushdown automaton over the input alphabet  $\Sigma\{a, b\}$  and stack alphabet  $\Gamma = \{\#, A\}$ .



The number of strings of length 100 accepted by the above pushdown automaton is \_\_\_\_\_.

**Ans. (50)**

The machine is pushing a “A” for every ‘a’ in input and popping a “A” for every ‘b’ in input which follows a’s. No a’s are allowed after the b’s and also at end of input, the stack must have at least one “A” for string to be accepted, which means number of a’s > number of b’s.

So language accepted is

$$L = \{a^m b^n \mid m > n \geq 0\}$$

Now strings of length 100 which satisfy above condition are  $\{a^{100}, a^{99}b, a^{98}b^2, \dots, a^{51}b^{48}\}$ .  
 Number of such strings is therefore 50.

**End of Solution**

**Q.9** Suppose that  $L_1$  is a regular language and  $L_2$  is a context-free language. Which one of the following languages is NOT necessarily context-free?

- (a)  $L_1 \cdot L_2$
- (b)  $L_1 \cup L_2$
- (c)  $L_1 - L_2$
- (d)  $L_1 \cap L_2$

**Ans. (c)**

$L_1 \rightarrow \text{Reg}$

$L_2 \rightarrow \text{Context-free}$

(a)  $L_1 \cdot L_2 = \text{Reg} \cdot \text{CFL} = \text{CFL} \cdot \text{CFL} = \text{CFL}$

(b)  $L_1 \cup L_2 = \text{Reg} \cup \text{CFL} = \text{CFL}$

(c)  $L_1 - L_2 = L_1 \cap \overline{L_2} = \text{Reg} \cap \overline{\text{CFL}} = \text{Reg} \cap \overline{\text{CSL}}$   
 $= \text{Reg} \cap \text{CSL} = \text{CSL}$

(d)  $L_1 \cap L_2 = \text{Reg} \cap \text{CFL} = \text{CFL}$

So (c) is not necessarily CFL.

**End of Solution**

**Q.10** A TCP server application is programmed to listen on port number  $P$  on host  $S$ . A TCP client is connected to the TCP server over the network.

Consider that while the TCP connection was active, the server machine  $S$  crashed and rebooted. Assume that the client does not use the TCP keepalive timer.

Which of the following behaviors is/are possible?

- (a) If the client sends a packet after the server reboot, it will receive a FIN segment.
- (b) If the client was waiting to receive a packet, it may wait indefinitely.
- (c) The TCP server application on  $S$  can listen on  $P$  after reboot.
- (d) If the client sends a packet after the server reboot, it will receive a RST segment.

**Ans. (b, c, d)**

**(a) False:** The situation resolves itself when client tries to send data to server over the dead connection, and server replies with an RST packet (not FIN), causing client to finally to close the connection forcibly.

FIN is used to close TCP connections gracefully in each direction (normal close of connection), while TCP RST is used in a scenario where TCP connections cannot recover from errors and the connection needs to reset forcibly.

**(b) True:** Since broken connections can only be detected by sending data, the receiving side will wait forever. This scenario is called a “half-open connection” because one side realizes the connection was lost but the other side believes it is still active.

**(c) True:** Yes, a TCP Server can listen to same port number even after reboot. For example, the SMTP service application usually listens on TCP port 25 for incoming requests. So, even after reboot the port 25 is assigned to SMTP.

**(d) True:** The situation resolves itself when client tries to send data to server over the dead connection, and server replies with an RST packet (not FIN).

**End of Solution**

**Q.11** There are 6 jobs with distinct difficulty levels, and 3 computers with distinct processing speeds. Each job is assigned to a computer such that:

- The fastest computer gets the toughest job and the slowest computer gets the easiest job.
- Every computer gets at least one job.

The number of ways in which this can be done is \_\_\_\_\_.

**Ans. (65)**

Let computers be A, B and C

Toughest job assigned to fastest computer (Say, A) is 1 way.

Easiest job assigned to shortest computer (Say, B) is 1 way.

Remaining 4 jobs to be assigned to 3 computers so that the computer C gets at least one job, since A and B already assigned a job.

Number of ways 4 jobs assigned to 3 computers =  $3^4$ .

Number of ways 4 jobs assigned to 3 computers, so that computer C does not get any job =  $2^4$ .

Required number of ways =  $3^4 - 2^4 = 81 - 16 = 65$  ways

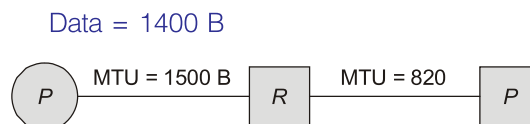
**End of Solution**

**Q.12** Consider two hosts  $P$  and  $Q$  connected through a router  $R$ . The maximum transfer unit (MTU) value of the link between  $P$  and  $R$  is 1500 bytes, and between  $R$  and  $Q$  is 820 bytes. A TCP segment of size 1400 bytes was transferred from  $P$  to  $Q$  through  $R$ , with IP identification value as 0x1234. Assume that the IP header size is 20 bytes. Further, the packet is allowed to be fragmented, i.e., Don't Fragment (DF) flag in the IP header is not set by  $P$ .

Which of the following statements is/are correct?

- (a) Two fragments are created at  $R$  and the IP datagram size carrying the second fragment is 620 bytes.
- (b) If the second fragment is lost,  $P$  is required to resend the whole TCP segment.
- (c) TCP destination port can be determined by analysing only the second fragment.
- (d) If the second fragment is lost,  $R$  will resend the fragment with the IP identification value 0x1234

**Ans. (a, b)**



For First Link – PR – There will be no fragment

For Second Link – RQ

$$\text{First Fragment} = 800(\text{Data}) + 20(\text{Header}) = 820 \text{ B}$$

$$\text{Second Fragment} = 1400 - 800 = 600 + 20(\text{H}) = 620 \text{ B}$$

**End of Solution**

**Q.13** For a Turing machine  $M$ ,  $\langle M \rangle$  denotes an encoding of  $M$ . Consider the following two languages:

$$L_1 = \{\langle M \rangle \mid M \text{ takes more than 2021 steps on all inputs}\}$$

$$L_2 = \{\langle M \rangle \mid M \text{ takes more than 2021 steps on some input}\}$$

Which one of the following options is correct?

- (a) Both  $L_1$  and  $L_2$  are decidable.
- (b)  $L_1$  is undecidable and  $L_2$  is decidable.
- (c) Both  $L_1$  and  $L_2$  are undecidable.
- (d)  $L_1$  is decidable and  $L_2$  is undecidable.

**Ans. (a)**

$$L_1 = \{\langle M \rangle \mid M \text{ takes more than 2021 steps on all inputs}\}$$

$$L_2 = \{\langle M \rangle \mid M \text{ takes more than 2021 steps on some input}\}$$

A Turing machine reads at most 2021 bits of input while making 2021 steps. So the halting behaviour is completely determined by the first 2021 bits of input. Now the number of strings with 2021 bits is finite and so generate all of them and in finite amount of time we can check if the given TM,  $M$  halts on any of these strings.

For  $L_1$ , the algorithm will be as follows,

Does not halts on all of these strings  $\rightarrow$  Yes

Halts on at least one of these strings  $\rightarrow$  No

For  $L_2$ , the algorithm will be as follows,  
does not halt on at least one of these strings  $\rightarrow$  Yes  
Halts on all of these strings  $\rightarrow$  No  
So, both  $L_1$  and  $L_2$  are decidable.

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**End of Solution**

**Q.14** In an undirected connected planar graph  $G$ , there are eight vertices and five faces. The number of edges in  $G$  is \_\_\_\_\_.

**Ans. (11)**

In a connected planar graph

$$r = e - n + 2$$

Here,  $n = 8, r = 5$

$$\therefore 5 = e - 8 + 2$$
$$e = 11$$

---

**End of Solution**

**Q.15** Consider the following array:

23	32	45	69	72	73	89	97
----	----	----	----	----	----	----	----

Which algorithm out of the following options uses the least number of comparisons (among the array elements) to sort the above array in ascending order?

- (a) Quicksort using the last element as pivot
- (b) Selection sort
- (c) Mergesort
- (d) Insertion sort

**Ans. (d)**

Given array already sorted.

$\therefore$  Insertion sort takes least number of comparisons  $\theta(n)$ .

Since, for a number which is to be inserted in the already sorted array, only 1 comparison will be required.

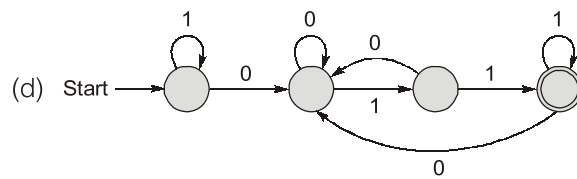
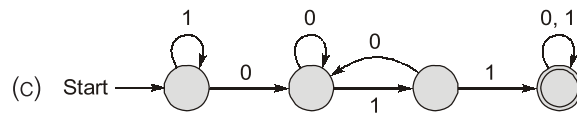
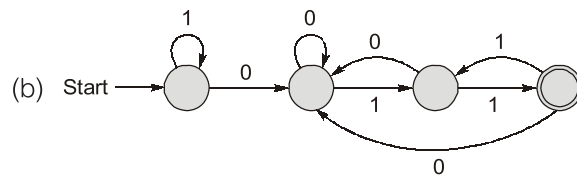
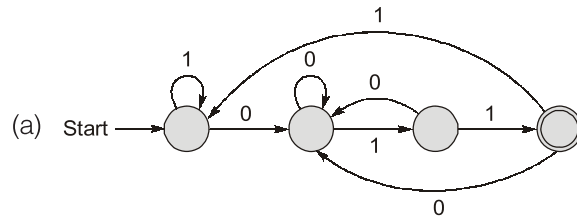
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**End of Solution**

**Q.16** Consider the following language:

$$L = \{w \in \{0,1\}^* \mid w \text{ ends with the substring } 011\}$$

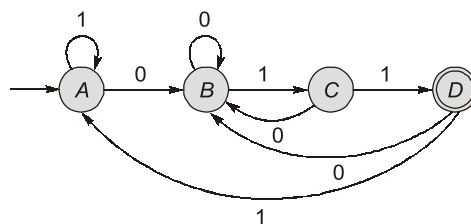
Which one of the following deterministic finite automata accepts  $L$ ?



**Ans. (a)**

Since 4 states are required in minimal DFA for this language “Ending with 011”, and since all given DFA's have 4 states, the answer must be same as minimal DFA.

The minimal DFA for this language is given below:



Which is same as option (a).

**End of Solution**

**Q.17** Let  $p$  and  $q$  be two propositions. Consider the following two formulae in propositional logic.

$$S_1 : (\neg p \wedge (p \vee q)) \rightarrow p$$

$$S_2 : q \rightarrow (\neg p \wedge (p \vee q))$$

Which one of the following choices is correct?

- (a) Neither  $S_1$  nor  $S_2$  is a tautology.
- (b) Both  $S_1$  and  $S_2$  are tautologies.
- (c)  $S_1$  is a tautology but  $S_2$  is not a tautology.
- (d)  $S_1$  is not a tautology but  $S_2$  is a tautology.

**Ans.** (c)

$$S_1 : (\neg p \wedge (p \vee q)) \rightarrow p$$

$$S_2 : q \rightarrow (\neg p \wedge (p \vee q))$$

$$S_1 : [p' (p + q)] \rightarrow p$$

$$\equiv (p'p + p'q) \rightarrow q$$

$$\equiv p'q \rightarrow q$$

$$\equiv (p'q)' + q$$

$$\equiv p + q' + q$$

$$\equiv p + 1$$

$$\equiv 1 \text{ (tautology)}$$

$$S_2 : q \rightarrow (p' (p + q))$$

$$\equiv q \rightarrow (p'p + p'q)$$

$$\equiv q \rightarrow p'q$$

$$\equiv q' + p'q$$

$$\equiv (q' + p') (q' + q)$$

$$\equiv q' + p' \text{ (contingency) (not a tautology)}$$

So, option (c)  $S_1$  is a tautology and  $S_2$  is not a tautology is correct.

---

**End of Solution**



**Q.18** Suppose a database system crashes again while recovering from a previous crash. Assume checkpointing is not done by the database either during the transactions or during recovery.

Which of the following statements is/are correct?

- (a) The same undo and redo list will be used while recovering again.
- (b) The system cannot recover any further.
- (c) The database will become inconsistent.
- (d) All the transactions that are already undone and redone will not be recovered again.

**Ans. (a)**

**End of Solution**

**Q.19** A relation  $R$  is said to be circular if  $aRb$  and  $bRc$  together imply  $cRa$ .

Which of the following options is/are correct?

- (a) If a relation  $S$  is reflexive and circular, then  $S$  is an equivalence relation.
- (b) If a relation  $S$  is circular and symmetric, then  $S$  is an equivalence relation.
- (c) If a relation  $S$  is transitive and circular, then  $S$  is an equivalence relation.
- (d) If a relation  $S$  is reflexive and symmetric, then  $S$  is an equivalence relation.

**Ans. (a)**

Let  $S$  be reflexive and circular,

Let us check symmetry:

**Symmetry:**

Let  $xSy$

Now since  $S$  is reflexive  $ySy$  is true.

So  $xSy$  and  $ySy$  is true

Now by circular property we get,  $ySx$

So  $xSy \Rightarrow ySx$

So  $S$  is symmetric.

**Transitive:**

Let  $xSy$  and  $ySz$

Now by circular property we get  $zSx$  and by symmetry property proved above, we get

$$zSx \Rightarrow xSz$$

So  $xSy$  and  $ySz \Rightarrow xSz$

So  $S$  is transitive.

So  $S$  is reflexive, symmetric and transitive and hence an equivalence relation.

So option (a) is true.

Option (b): Let  $S$  be circular and symmetric.

Let  $S$  be defined on set  $\{1, 2, 3\}$

Now empty relation is circular and symmetric but not reflexive. So  $S$  need not be an equivalence relation.

So option (b) is false.

Option (c): Let  $S$  be transitive and circular.

Let  $S$  be defined on the set  $\{1, 2, 3\}$

Now empty relation again satisfies transitive and circular but is not reflexive. So  $S$  need not be an equivalence relation.

So option (c) is false.

Option (d): Reflexive and symmetric need not be transitive for example on  $\{1, 2, 3\}$ .

$$S = \{(1, 1), (2, 2), (3, 3), (1, 2), (2, 1), (2, 3), (3, 2)\}$$

is reflexive and symmetric. But it is not transitive because  $(1, 2)$  and  $(2, 3)$  belong to  $S$  but  $(1, 3)$  does not.

So option (d) is false.

**End of Solution**

**Q.20** Let  $P$  be an array containing  $n$  integers. Let  $t$  be the lowest upper bound on the number of comparisons of the array elements, required to find the minimum and maximum values in an arbitrary array of  $n$  elements. Which one of the following choices is correct?

(a)  $t > n$  and  $t \leq 3\left\lceil \frac{n}{2} \right\rceil$

(b)  $t > 3\left\lceil \frac{n}{2} \right\rceil$  and  $t \leq 2n - 2$

(c)  $t > 2n - 2$

(d)  $t > \lceil \log_2(n) \rceil$  and  $t \leq n$

**Ans. (b)**

$$t > 3\left\lceil \frac{n}{2} \right\rceil \text{ and } t \leq 2n - 2$$

Using straight max-min algo, WC number of comparisons =  $2n - 2$ .

Using divide and conquer min-max algo, WC number of comparisons =  $\frac{3n}{2} - 2$ .

**[Note: But the official answer key given by IIT-Bombay is option (a)]**

**End of Solution**

**Q.21** In the context of operating systems, which of the following statements is/are correct with respect to paging?

(a) Paging incurs memory overheads.

(b) Multi-level paging is necessary to support pages of different sizes.

(c) Page size has no impact on internal fragmentation.

(d) Paging helps solve the issue of external fragmentation.

**Ans. (a, d)**

**End of Solution**

**Q.22** The lifetime of a component of a certain type is a random variable whose probability density function is exponentially distributed with parameter 2. For a randomly picked component of this type, the probability that its lifetime exceeds the expected lifetime (rounded to 2 decimal places) is \_\_\_\_\_.

**Ans. (0.367)**

Let,  $t = \{\text{lifetime of component}\}$  and  $\mu = 2$

$$\text{Expected lifetime} = \frac{1}{\mu} \text{ and } f(t) = \begin{cases} \mu e^{-\mu t}, & t > 0 \\ 0, & \text{otherwise} \end{cases}$$

$$\begin{aligned} P\left(t > \frac{1}{\mu}\right) &= \int_{\frac{1}{\mu}}^{\infty} f(t) dt = \int_{\frac{1}{\mu}}^{\infty} \mu e^{-\mu t} dt = -\left[e^{-\mu t}\right]_{\frac{1}{\mu}}^{\infty} = -\left[0 - e^{-1}\right] \\ &= \frac{1}{e} = \frac{1}{2.713} = 0.367 \end{aligned}$$

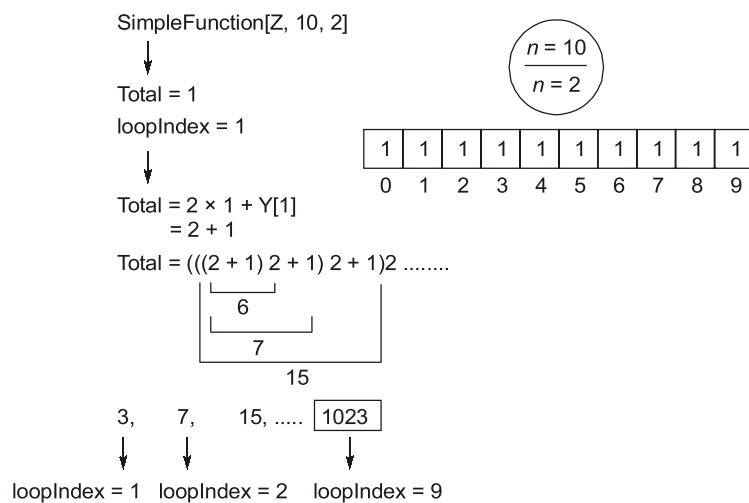
**End of Solution**

**Q.23** Consider the following ANSI C function:

```
int SimpleFunction (int Y[ ], int n, int x)
{
    int total = Y[0], loopIndex;
    for (loopIndex = 1; loopIndex <= n - 1; loopIndex++)
        total = x * total + Y[loopIndex];
    return total;
}
```

Let Z be an array of 10 elements with  $Z[i] = 1$ , for all  $i$  such that  $0 \leq i \leq 9$ . The value returned by SimpleFunction (Z, 10, 2) is \_\_\_\_\_.

**Ans. (1023)**



**End of Solution**

- Q.24** Consider a linear list based directory implementation in a file system. Each directory is a list of nodes, where each node contains the file name along with the file metadata, such as the list of pointers to the data blocks. Consider a given directory foo.
- Which of the following operations will necessarily require a full scan of foo for successful completion?
- (a) Opening of an existing file in foo    (b) Renaming of an existing file in foo  
 (c) Creation of a new file in foo        (d) Deletion of an existing file from foo

**Ans. (a, b, d)**

[Note: But the official answer key given by IIT-Bombay is option (b, c)]

**End of Solution**

- Q.25** Let  $G = (V, E)$  be an undirected unweighted connected graph. The diameter of  $G$  is defined as

$$\text{diam}(G) = \max_{u, v \in V} \{\text{the length of shortest path between } u \text{ and } v\}$$

Let  $M$  be the adjacency matrix of  $G$ .

Define graph  $G_2$  on the same set of vertices with adjacency matrix  $N$ , where

$$N_{ij} = \begin{cases} 1 & \text{if } M_{ij} > 0 \text{ or } P_{ij} > 0, \text{ where } P = M^2 \\ 0 & \text{Otherwise} \end{cases}$$

Which one of the following statements is true?

- (a)  $\lceil \text{diam}(G)/2 \rceil < \text{diam}(G_2) < \text{diam}(G)$   
 (b)  $\text{diam}(G_2) \leq \lceil \text{diam}(G)/2 \rceil$   
 (c)  $\text{diam}(G) < \text{diam}(G_2) \leq 2 \text{diam}(G)$   
 (d)  $\text{diam}(G_2) = \text{diam}(G)$

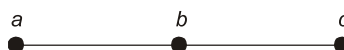
**Ans. (b)**

$$N_{ij} = \begin{cases} 1 & \text{if } M_{ij} > 0 \text{ or } P_{ij} > 0, \text{ where } P = M^2 \\ 0 & \end{cases}$$

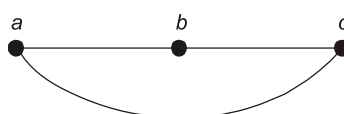
Means  $G_2$  will have not only all the edges of in  $G$ , but also will have edges connecting vertices in  $G$  which have a path of length 2, since  $M^2$  will have all edges between  $u$  and  $v$  if there is a path of length 2 between shown in  $G$ .

**Option (a)**  $\lceil \text{diam}(G)/2 \rceil < \text{diam}(G_2) < \text{diam}(G)$

Let  $G$  be the graph shown below with  $\text{dia}(G) = 2$



Now  $G_2$  will be



Since  $(a, c)$  has a path of length 2 in  $G$ ,  $G_2$  will have an edge connecting  $a$  and  $c$ .

Now diameter of  $G_2 = 1$

This violates option (a) condition that  $\text{dia}(G_2) > \text{dia}(G)/2$ .

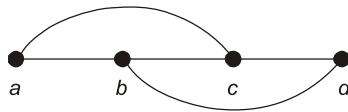
So option (a) false.

**Option (b)**  $\text{dia}(G_2) \leq \lceil \text{dia}(G)/2 \rceil$

Consider graph below with  $\text{dia}(G) = 3$



Now  $G_2$  will connect all paths of length 2 with edges and will be



Now  $\text{dia}(G_2) = 2$

But  $\text{dia}(G) = 3$

$$\text{dia}(G_2) \leq \lceil \text{dia}(G)/2 \rceil$$

$$2 \leq \lceil 1.5 \rceil$$

$$2 \leq 2$$

Satisfied.

Option (b) is correct.

**Option (c)**  $\text{dia}(G) < \text{dia}(G_2) \leq 2 \text{ dia}(G)$

Taking previous option example where  $\text{dia}(G) = 3$  and  $\text{dia}(G_2) = 2$

$$3 \leq 2 \leq 6$$

Is false.

So option (c) is false.

**Option (d)**  $\text{diam}(G_2) = \text{diam}(G)$

Taking previous example where  $\text{dia}(G) = 3$  and  $\text{dia}(G_2) = 2$

$$2 = 3 \text{ is false}$$

So option (d) is false.

So correct option is option (b).

---

End of Solution

**Q.26** Consider the following context-free grammar where the set of terminals is  $\{a, b, c, d, f\}$

$$S \rightarrow daT \mid Rf$$

$$T \rightarrow aS \mid baT \mid \epsilon$$

$$R \rightarrow caTR \mid \epsilon$$

The following is a partially-filled LL(1) parsing table.

	a	b	c	d	f	\$
S			①	$S \rightarrow daT$	②	
T	$T \rightarrow aS$	$T \rightarrow baT$	③		$T \rightarrow \epsilon$	④
R			$R \rightarrow caTR$		$R \rightarrow \epsilon$	

Which one of the following choices represents the correct combination for the numbered cells in the parsing table ("blank" denotes that the corresponding cell is empty)?

- (a) ①  $S \rightarrow Rf$     ② blank    ③ blank    ④  $T \rightarrow \epsilon$
- (b) ① blank    ②  $S \rightarrow Rf$     ③  $T \rightarrow \epsilon$     ④  $T \rightarrow \epsilon$
- (c) ① blank    ②  $S \rightarrow Rf$     ③ blank    ④ blank
- (d) ①  $S \rightarrow Rf$     ②  $S \rightarrow Rf$     ③  $T \rightarrow \epsilon$     ④  $T \rightarrow \epsilon$

**Ans. (d)**

$$S \rightarrow daT \mid Rf$$

$$T \rightarrow aS \mid baT \mid \epsilon$$

$$R \rightarrow caTR \mid \epsilon$$

$$\text{First}(S) = \{d, c, f\}$$

$$\text{Follow}(S) = \{c, f, \$\}$$

$$\text{First}(T) = \{a, b, \epsilon\}$$

$$\text{Follow}(T) = \{c, f, \$\}$$

$$\text{First}(R) = \{c, \epsilon\}$$

$$\text{Follow}(R) = \{f\}$$

	a	b	c	d	f	\$
S			① $S \rightarrow Rf$	$S \rightarrow daT$	② $S \rightarrow Rf$	
T	$T \rightarrow aS$	$T \rightarrow baT$	③ $T \rightarrow \epsilon$		$T \rightarrow \epsilon$	④ $T \rightarrow \epsilon$
R			$R \rightarrow caTR$		$R \rightarrow \epsilon$	

- ①  $S \rightarrow Rf$     ②  $S \rightarrow Rf$     ③  $T \rightarrow \epsilon$     ④  $T \rightarrow \epsilon$

Hence option (d) is answer.

**End of Solution**

**Q.27** Let  $\langle M \rangle$  denote an encoding of an automaton  $M$ . Suppose that  $\Sigma = \{0, 1\}$ . Which of the following languages is/are NOT recursive?

- (a)  $L = \{\langle M \rangle \mid M \text{ is a DFA such that } L(M) = \Sigma^*\}$
- (b)  $L = \{\langle M \rangle \mid M \text{ is a PDA such that } L(M) = \Sigma^*\}$
- (c)  $L = \{\langle M \rangle \mid M \text{ is a PDA such that } L(M) = \emptyset\}$
- (d)  $L = \{\langle M \rangle \mid M \text{ is a DFA such that } L(M) = \emptyset\}$

**Ans. (b)**

- Option (a): Completeness problem for regular  $\rightarrow$  Decidable  
 $\rightarrow$  Recursive
- Option (b): Completeness problem for CFL's  $\rightarrow$  Undecidable  
 $\rightarrow$  Not recursive
- Option (c): Emptiness problem for CFL's  $\rightarrow$  Decidable  
 $\rightarrow$  Recursive
- Option (d): Emptiness problem for Regular  $\rightarrow$  Decidable  
 $\rightarrow$  Recursive

So only option (b) is not recursive.

**End of Solution**

**Q.28** Consider the following instruction sequence where registers R1, R2 and R3 are general purpose and MEMORY[X] denotes the content at the memory location X.

Instruction	Semantics	Instruction Size (bytes)
MOV R1, (5000)	$R1 \leftarrow \text{MEMORY}[5000]$	4
MOV R2, (R3)	$R2 \leftarrow \text{MEMORY}[R3]$	4
ADD R2, R1	$R2 \leftarrow R1 + R2$	2
MOV (R3), R2	$\text{MEMORY}[R3] \leftarrow R2$	4
INC R3	$R3 \leftarrow R3 + 1$	2
DEC R1	$R1 \leftarrow R1 - 1$	2
BNZ 1004	Branch if not zero to the given absolute address	2
HALT	Stop	1

Assume that the content of the memory location 5000 is 10 and the content of the register R3 is 3000. The content of each of the memory locations from 3000 to 3010 is 50. The instruction sequence starts from the memory location 1000. All the numbers are in decimal format. Assume that the memory is byte addressable.

After the execution of the program, the content of memory location 3010 is \_\_\_\_\_.

**Ans. (50)**

Program execution 10 memory cells information is accessed for reading & writing starting from 3000 location.

So, 3000 to 3009 cells are accessed for read and write.

$\therefore$  No change in [3010] cell.

So, it contain 50 only.

**End of Solution**

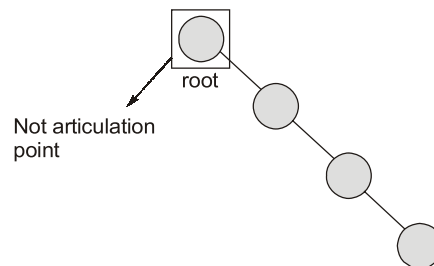
**Q.29** An *articulation point* in a connected graph is a vertex such that removing the vertex and its incident edges disconnects the graph into two or more connected components. Let  $T$  be a DFS tree obtained by doing DFS in a connected undirected graph  $G$ .

Which of the following options is/are correct?

- (a) Root of  $T$  is an articulation point in  $G$  if and only if it has 2 or more children.
- (b) A leaf of  $T$  can be an articulation point in  $G$ .
- (c) Root of  $T$  can never be an articulation point in  $G$ .
- (d) If  $u$  is an articulation point in  $G$  such that  $x$  is an ancestor of  $u$  in  $T$  and  $y$  is a descendent of  $u$  in  $T$ , then all paths from  $x$  to  $y$  in  $G$  must pass through  $u$ .

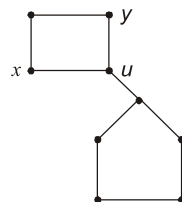
**Ans. (a)**

- (a) True



We need at least 2 children so that root is articulation point.

- (b) False: This can never happen. Leaf will always have degree = 1.
- (c) False: Check option (a) for more information.
- (d) False: Below is the reasoning to show how this is false.



If  $u$  is articulation point, then removing  $u$  generates 2 connected components, now there might be a case when  $x$  and  $y$  will belong to either one of the connected component and hence a path will exist between them without passing through  $u$ .

Option (a) is correct.

**End of Solution**



**Q.30** Consider the relation  $R(P, Q, S, T, X, Y, Z, W)$  with the following functional dependencies

$$PQ \rightarrow X; P \rightarrow YX; Q \rightarrow Y; Y \rightarrow ZW$$

Consider the decomposition of the relation  $R$  into the constituent relations according to the following two decomposition schemes.

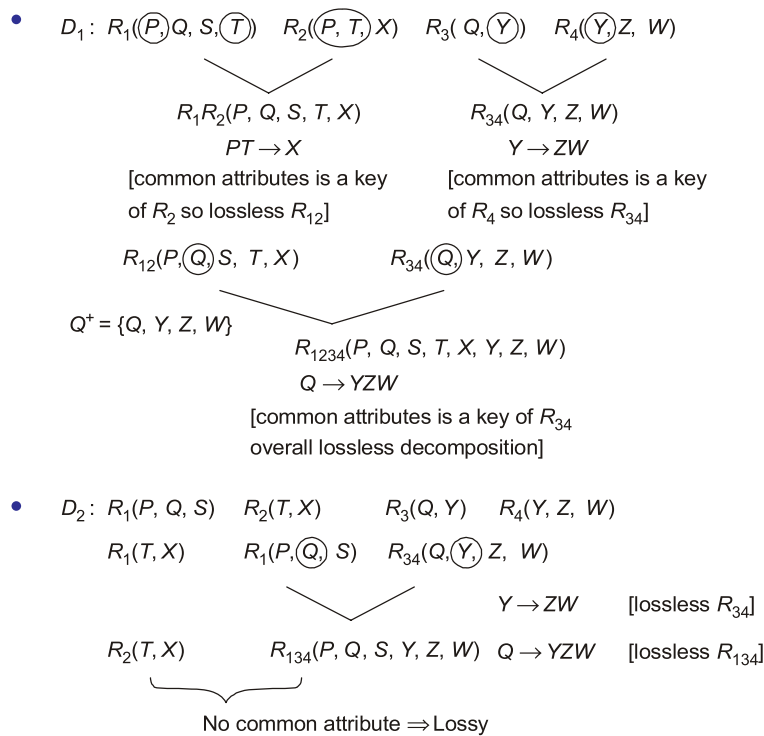
$$D_1 : R = [(P, Q, S, T); (P, T, X); (Q, Y); (Y, Z, W)]$$

$$D_2 : R = [(P, Q, S); (T, X); (Q, Y); (Y, Z, W)]$$

Which one of the following options is correct?

- (a)  $D_1$  is a lossy decomposition, but  $D_2$  is a lossless decomposition.
- (b) Both  $D_1$  and  $D_2$  are lossless decompositions.
- (c) Both  $D_1$  and  $D_2$  are lossy decompositions.
- (d)  $D_1$  is a lossless decomposition, but  $D_2$  is a lossy decomposition.

**Ans. (d)**



$D_1$  is a lossless decomposition, but  $D_2$  is a lossy decomposition.

End of Solution

**Q.31** A five-stage pipeline has stage delays of 150, 120, 150, 160 and 140 nanoseconds. The registers that are used between the pipeline stages have a delay of 5 nanoseconds each.

The total time to execute 100 independent instructions on this pipeline, assuming there are no pipeline stalls, is \_\_\_\_\_ nanoseconds.

**Ans. (17160)**

$$\begin{aligned}
 K &= 5 \\
 t_p &= \max(\text{Stage delay} + \text{Buffer delay}) = 165 \text{ ns} \\
 n &= 100 \text{ (finite)} \\
 &\downarrow \\
 ET_{\text{pipe}} &= (K + n - 1)t_p \\
 &= (5 + 100 - 1) 165 \text{ ns} = 17160 \text{ ns}
 \end{aligned}$$

**End of Solution**

**Q.32** Consider the sliding window flow-control protocol operating between a sender and a receiver over a full-duplex error-free link. Assume the following:

- The time taken for processing the data frame by the receiver is negligible.
- The time taken for processing the acknowledgement frame by the sender is negligible.
- The sender has infinite number of frames available for transmission.
- The size of the data frame is 2000 bits and the size of the acknowledgement frame is 10 bits.
- The link data rate in each direction is 1 Mbps (=  $10^6$  bits per second).
- One way propagation delay of the link is 100 milliseconds.

The minimum value of the sender's window size in terms of the number of frames, (rounded to the nearest integer) needed to achieve a link utilization of 50% is \_\_\_\_\_.

**Ans. (51)**

$$\begin{aligned}
 T_t(\text{packet}) &= \frac{L}{B.W} \\
 \Rightarrow \frac{2000 \text{ bits}}{10^6 \text{ bps}} &= 2 \times 10^{-3} \text{ sec} = 2 \text{ millisecc} \\
 T_t(\text{Ack}) &= \frac{L}{B.W} \\
 \Rightarrow \frac{10 \text{ bits}}{10^6 \text{ bps}} &= 10^{-5} \text{ sec} = 10^{-2} \text{ millisecc} = 0.01 \text{ millisecc} \\
 T_p &= 100 \text{ millisecc} \\
 \text{Total time} &= T_t(\text{packet}) + T_p + T_t(\text{Ack}) \\
 &= 2 + 100 + 0.01 = 102.01 \text{ millisecc} \\
 \text{Efficiency} &= 50\% = \frac{1}{2} \\
 \text{Efficiency} &= \frac{\text{Useful time}}{\text{Total time}}
 \end{aligned}$$

$$\frac{1}{2} = \frac{n \times T_t}{\text{Total time}}$$

$$\Rightarrow 2 \times n = 102.01$$

$$\Rightarrow n = \frac{102.01}{2} \Rightarrow 51.005$$

For minimum we have to take ceil, hence size of window = 51.

**End of Solution**

**Q.33** Consider the following sequence of operations on an empty stack.

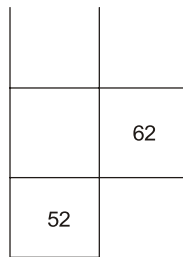
push(54); push(52); pop( ); push(55); push(62); s = pop( ):

Consider the following sequence of operations on an empty queue.

enqueue(21); enqueue(24); dequeue( ); enqueue(28); enqueue(32); q = dequeue( ):

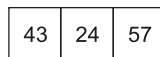
The value of s + q is \_\_\_\_\_.

**Ans. (86)**



$$S = 62$$

$$R = 24$$



$$S + R = 86$$

**End of Solution**

**Q.34** Let the representation of a number in base 3 be 210. What is the hexadecimal representation of the number?

(a) 528

(b) 21

(c) 15

(d) D2

**Ans. (c)**

$$\begin{aligned} (210)_3 &= 3^2 \times 2 + 3^1 \times 1 + 3^0 \times 0 \\ &= 18 + 3 + 0 \\ &= (21)_{10} \end{aligned}$$

$$(21)_{10} = 16 \left| \begin{array}{l} 21 \\ 1 \rightarrow 5 \end{array} \right. = (15)_{16}$$

**End of Solution**

**Q.35** A relation  $r(A, B)$  in a relational database has 1200 tuples. The attribute  $A$  has integer values ranging from 6 to 20, and the attribute  $B$  has integer values ranging from 1 to 20. Assume that the attributes  $A$  and  $B$  are independently distributed. The estimated number of tuples in the output of  $\sigma_{(A > 10) \vee (B = 18)}(r)$  is \_\_\_\_\_.

**Ans. (820)**

- There are 10 distinct integer for  $(A > 10)$  out of 15.
- There are 1 distinct integer for  $(B = 18)$  out of 20.

$$P(A > 10) = \frac{10}{15} = \frac{2}{3}$$

$$P(B = 18) = \frac{1}{20}$$

$$P(A > 10) \wedge (B = 18) = \frac{2}{3} \times \frac{1}{20} = \frac{1}{30}$$

Now,

$$\begin{aligned} P((A > 10) \vee (B = 18)) &= P(A > 10) + P(B = 18) - P(A > 10) \wedge (B = 18) \\ &= \frac{2}{3} + \frac{1}{20} - \frac{1}{30} = \frac{40+3-2}{60} = \frac{41}{60} \end{aligned}$$

$$\text{Estimated number of tuples} = \frac{41}{60} \times 1200 = 820 \text{ tuples}$$

**End of Solution**

**Q.36** Assume that a 12-bit Hamming codeword consisting of 8-bit data and 4 check bits is  $d_8 d_7 d_6 d_5 c_8 d_4 d_3 d_2 c_4 d_1 c_2 c_1$ , where the data bits and the check bits are given in the following tables:

Data bits							
$d_8$	$d_7$	$d_6$	$d_5$	$d_4$	$d_3$	$d_2$	$d_1$
1	1	0	$x$	0	1	0	1

Check bits			
$c_8$	$c_4$	$c_2$	$c_1$
$y$	0	1	0

Which one of the following choices gives the correct values of  $x$  and  $y$ ?

- (a)  $x$  is 0 and  $y$  is 0                      (b)  $x$  is 0 and  $y$  is 1  
 (c)  $x$  is 1 and  $y$  is 0                      (d)  $x$  is 1 and  $y$  is 1

Ans. (a)

1	2	3	4	5	6	7	8		$c_1$	$c_2$	$c_4$	$c_8$
1	0	1	0	$x$	0	1	1		0	1	0	$y$

$c_1 =$	1	3	5	7	9	11						
	0	1	0	0	$x$	1						

$\leftarrow x=0 \rightarrow$

1	2	3	4	5	6	7	8	9	10	11	12	
0	1	1	0	0	1	0	$y$	$x$	0	1	1	

$c_8 =$	8	9	10	11	12							
	0	0	0	1	1							

$\leftarrow y=0 \rightarrow$

End of Solution

**Q.37** Let  $r_i(z)$  and  $w_i(z)$  denote read and write operations respectively on a data item  $z$  by a transaction  $T_i$ . Consider the following two schedules:

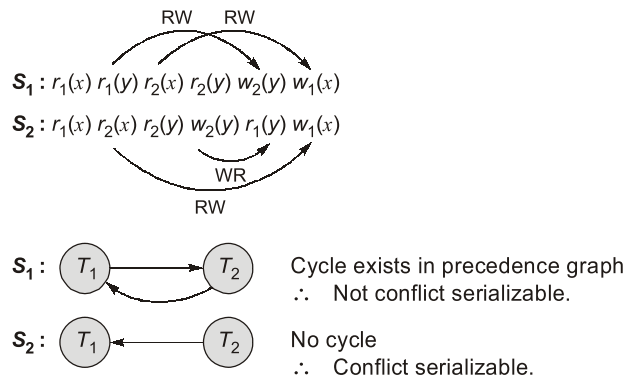
$S_1 : r_1(x), r_1(y), r_2(x), r_2(y), w_2(y), w_1(x)$

$S_2 : r_1(x), r_2(x), r_2(y), w_2(y), r_1(y), w_1(x)$

Which one of the following options is correct?

- (a)  $S_1$  is not conflict serializable, and  $S_2$  is conflict serializable.
- (b) Neither  $S_1$  nor  $S_2$  is conflict serializable.
- (c) Both  $S_1$  and  $S_2$  are conflict serializable.
- (d)  $S_1$  is conflict serializable, and  $S_2$  is not conflict serializable.

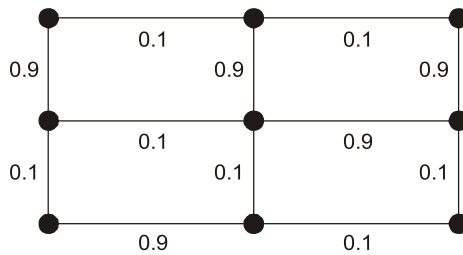
Ans. (a)



$S_1$  is not conflict serializable, and  $S_2$  is conflict serializable.

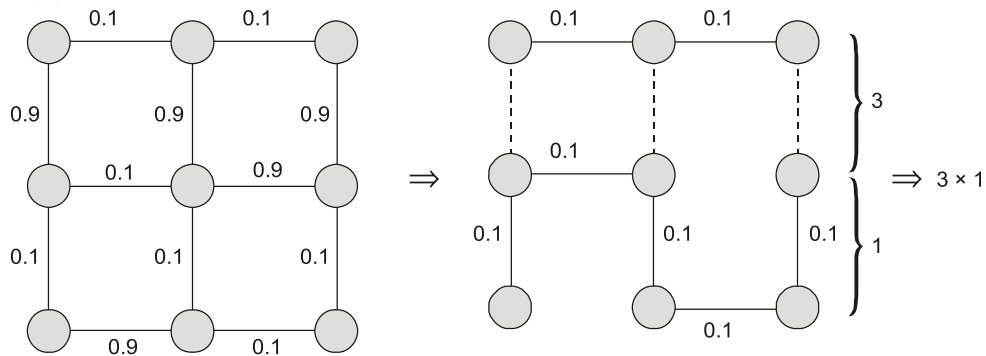
End of Solution

**Q.38** Consider the following undirected graph with edge weights as shown:



The number of minimum-weight spanning trees of the graph is \_\_\_\_\_.

**Ans. (3)**

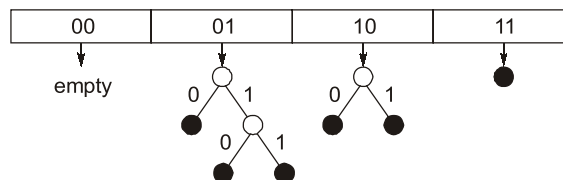


**End of Solution**

**Q.39** Consider a dynamic hashing approach for 4-bit integer keys:

1. There is a main hash table of size 4.
2. The 2 least significant bits of a key is used to index into the main hash table,
3. Initially, the main hash table entries are empty.
4. Thereafter, when more keys are hashed into it, to resolve collisions, the set of all keys corresponding to a main hash table entry is organized as a binary tree that grows on demand.
5. First, the 3<sup>rd</sup> least significant bit is used to divide the keys into left and right subtrees.
6. To resolve more collisions, each node of the binary tree is further sub-divided into left and right subtrees based on the 4<sup>th</sup> least significant bit.
7. A split is done only if it is needed, i.e., only when there is a collision.

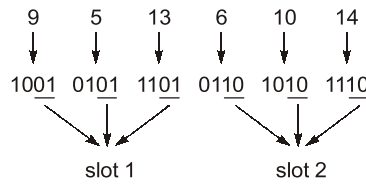
Consider the following state of the hash table.



Which of the following sequences of key insertions can cause the above state of the hash table (assume the keys are in decimal notation)?

- |                         |                        |
|-------------------------|------------------------|
| (a) 9, 5, 13, 6, 10, 14 | (b) 9, 5, 10, 6, 7, 1  |
| (c) 10, 9, 6, 7, 5, 13  | (d) 5, 9, 4, 13, 10, 7 |

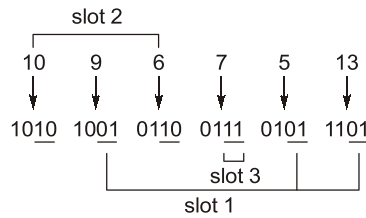
Ans. (c)



No entry for 11 in this option and 10 → has 3 entries but 2 are required. So option (a) is wrong.

Similarly, (b), (d) options can also be proved wrong. Option (c) is correct.

Let's check option (c).



End of Solution

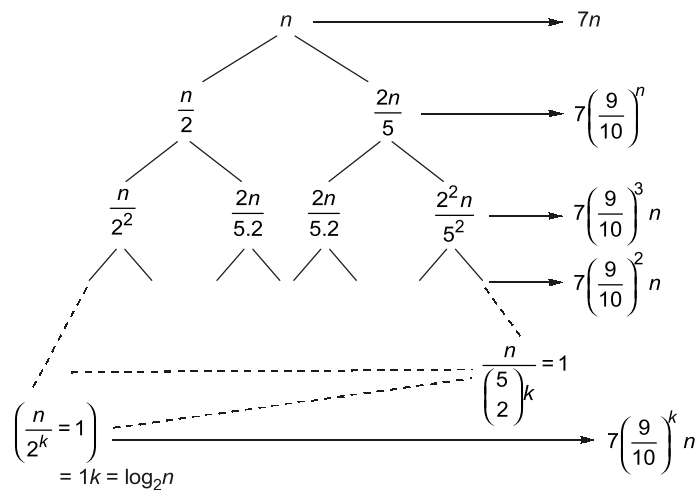
Q.40 Consider the following recurrence relation:

$$T(n) = \begin{cases} T\left(\frac{n}{2}\right) + T\left(\frac{2n}{5}\right) + 7n & \text{if } n > 0 \\ 1 & \text{if } n = 0 \end{cases}$$

Which one of the following options is correct?

- (a)  $T(n) = \theta(n \log n)$
- (b)  $T(n) = \theta(n^{5/2})$
- (c)  $T(n) = \theta(n)$
- (d)  $T(n) = \theta((\log n)^{5/2})$

Ans. (c)



$$T(n) = 7n \left[ 1 + \frac{9}{10} + \left(\frac{9}{10}\right)^2 + \left(\frac{9}{10}\right)^3 + \dots + \left(\frac{9}{10}\right)^k \right]$$

$$= \frac{7n \left[ 1 - \left(\frac{9}{10}\right)^{k+1} \right]}{1 - \frac{9}{10}} = \frac{7n}{\frac{1}{10}} \left[ 1 - \left(\frac{9}{10}\right)^{\log_2 n + 1} \right]$$

Decreasing term

$$\therefore T(n) = \theta(n)$$

**End of Solution**

- Q.41** Consider a computer system with a byte-addressable primary memory of size  $2^{32}$  bytes. Assume the computer system has a direct-mapped cache of size 32 KB (1 KB =  $2^{10}$  bytes), and each cache block is of size 64 bytes. The size of the tag field is \_\_\_\_\_ bits.

**Ans. (17)**

$$\text{MM size} = 2^{32} \text{ B}$$

$$\text{Addressable size} = \log_2 2^{32} = 32 \text{ bit}$$

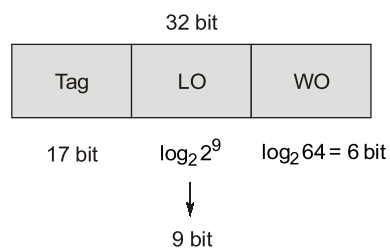
Direct mapped

$$\text{CM size} = 32 \text{ KB}$$

$$\text{Block size} = 64 \text{ B}$$

$$\therefore \text{Number of lines} = \frac{32 \text{ K}}{64} \Rightarrow \frac{2^{15}}{2^6} = 2^9$$

Address format:



**End of Solution**

- Q.42** Let  $G$  be a group of order 6, and  $H$  be a subgroup of  $G$  such that  $1 < |H| < 6$ . Which one of the following options is correct?
- Both  $G$  and  $H$  are always cyclic.
  - $G$  is always cyclic, but  $H$  may not be cyclic.
  - $G$  may not be cyclic, but  $H$  is always cyclic.
  - Both  $G$  and  $H$  may not be cyclic.



Ans. (c)

$$|G| = 6$$

$H$  is subgroup, so by Lagrange's theorem

$$|H| = 1, 2, 3 \text{ or } 6 \text{ (Divisor's of } 6)$$

Now it is given that  $1 < |H| < 6$

or  $|H| = 2 \text{ or } 3$

Since 2 and 3 are both prime and since every group of prime order is cyclic,  $H$  is surely cyclic.

But order of  $|G| = 6$  which is not prime.

So  $G$  may or may not be cyclic.

So  $G$  may not be cyclic, but  $H$  is always cyclic.

Option (c) is correct.

End of Solution

**Q.43** The following relation records the age of 500 employees of a company, where empNo (indicating the employee number) is the key:

empAge(empNo, age)

Consider the following relational algebra expression:

$$\Pi_{\text{empNo}} (\text{empAge} \bowtie_{(\text{age} > \text{age1})} \rho_{\text{empNo1, age1}} (\text{empAge}))$$

What does the above expression generate?

- (a) Employee numbers of all employees whose age is not the minimum.
- (b) Employee numbers of all employees whose age is the minimum.
- (c) Employee numbers of only those employees whose age is the maximum.
- (d) Employee numbers of only those employees whose age is more than the age of exactly one other employee.

Ans. (a)

empAge(empNo, age)

$$\Pi_{\text{empNo}} (\text{empAge} \bowtie_{(\text{age} > \text{age1})} \rho_{\text{empNo1, age1}} (\text{empAge}))$$

Retrieve empNo values of empAge those are having 'age' greater than some age.

↓

Retrieve Employee Number (empNo) of all employee whose age is greater than some employee' age.

↓

Employee number of all employees whose age is not the minimum. [Since it is greater than atleast 1 age].

End of Solution

**Q.44** Consider the following expression:

$$\lim_{x \rightarrow -3} \frac{\sqrt{2x+22} - 4}{x+3}$$

The value of the above expression (rounded to 2 decimal places) is \_\_\_\_\_.

**Ans. (0.25)**

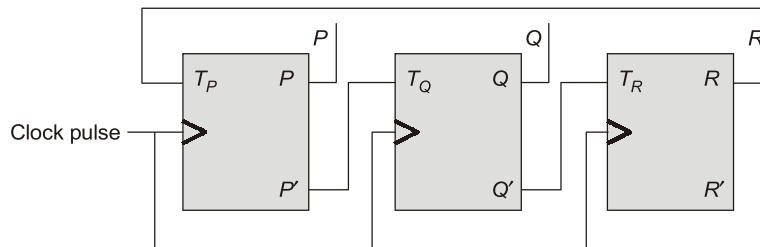
$$\lim_{x \rightarrow -3} \left( \frac{\sqrt{2x+22} - 4}{x+3} \right) \left( \frac{0}{0} \text{ form} \right), \text{ so apply L'Hospital's rule}$$

$$= \lim_{x \rightarrow -3} \left[ \frac{2}{2\sqrt{2x+22}} \right] = \lim_{x \rightarrow -3} \frac{1}{\sqrt{2x+22}}$$

$$= \frac{1}{\sqrt{16}} = \frac{1}{4} = 0.25$$

**End of Solution**

**Q.45** Consider a 3-bit counter, designed using T flip-flops, as shown below:



Assuming the initial state of the counter given by  $PQR$  as 000. What are the next three states?

- (a) 011, 101, 111                      (b) 001, 010, 000  
 (c) 001, 010, 111                      (d) 011, 101, 000

**Ans. (d)**

Clock	Present state			Flip-flop inputs			Next state		
	$P$	$Q$	$R$	$T_P = R$	$T_Q = \bar{P}$	$T_R = \bar{Q}$	$P^+$	$Q^+$	$R^+$
1	0	0	0	0	1	1	0	1	1
2	0	1	1	1	1	0	1	0	1
3	1	0	1	1	0	1	0	0	0

$\therefore$  Next three states of counter are 011, 101, 000  
 Hence option (d) is answer.

**End of Solution**

**Q.46** Consider the following statements:

$S_1$  : Every SLR(1) grammar is unambiguous but there are certain unambiguous grammars that are not SLR(1).

$S_2$  : For any context-free grammar, there is a parser that takes at most  $O(n^3)$  time to parse a string of length  $n$ .

Which one of the following options is correct?

- (a)  $S_1$  is true and  $S_2$  is true                      (b)  $S_1$  is false and  $S_2$  is false  
 (c)  $S_1$  is true and  $S_2$  is false                      (d)  $S_1$  is false and  $S_2$  is true

**Ans. (a)**

$S_1$  : Every SLR(1) is unambiguous but every unambiguous in not SLR(1). So  $S_1$  is true.

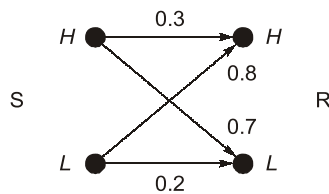
$S_2$  : Using CYK algorithm which takes at most  $O(n^3)$  time to parse a string, where  $n$  is the length of the string. So  $S_2$  is true.

Both  $S_1$  is true and  $S_2$  is true

**End of Solution**

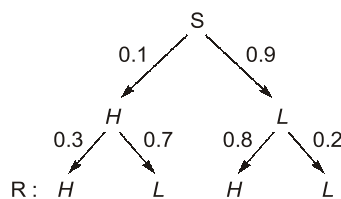
**Q.47** A sender (S) transmits a signal, which can be one of the two kinds:  $H$  and  $L$  with probabilities 0.1 and 0.9 respectively, to a receiver (R).

In the graph below, the weight of edge  $(u, v)$  is the probability of receiving  $v$  when  $u$  is transmitted, where  $u, v \in \{H, L\}$ . For example, the probability that the received signal is  $L$  given the transmitted signal was  $H$ , is 0.7.



If the received signal is  $H$ , the probability that the transmitted signal was  $H$  (rounded to 2 decimal places) is \_\_\_\_\_.

**Ans. (0.04)**



$$P\left(\frac{H_S}{H_R}\right) = \frac{P(H_S \cap H_R)}{P(H_R)}$$

$$= \frac{0.1 \times 0.3}{0.1 \times 0.3 + 0.9 \times 0.8} = \frac{1}{25} = 0.04$$

**End of Solution**

**Q.48** Consider the following ANSI C program:

```
#include <stdio.h>
int main( )
{
    int i, j, count;
    count = 0;
    i = 0;
    for (j = -3; j <= 3; j++)
    {
        if ((j >= 0) && (i++))
            count = count + j;
    }
    count = count + i;
    printf("%d", count);
    return 0;
}
```

Which one of the following options is correct?

- (a) The program will compile successfully and output 8 when executed.
- (b) The program will compile successfully and output 13 when executed.
- (c) The program will not compile successfully.
- (d) The program will compile successfully and output 10 when executed

**Ans. (d)**

$j = -3, j = -2, j = -1 \Rightarrow$  Short circuiting in  $\text{if}((j \geq 0) \ \&\& \ (i++))$

$j = 0; i = 0$  used but  $i$  is 1

$j = 1; i = 1$  used but  $i$  is 2

$\text{count} = 0 + 1 = 1$

$j = 2; i = 2$  used but  $i$  is 3

$\text{count} = 1 + 2$

$j = 3; i = 3$  used but  $i$  is 4

$\text{count} = 3 + 3 = 6$

$\text{count} = 6 + 4 = 10$  corresponding to  $\text{count} = \text{count} + i$

Option (d) is correct.

---

**End of Solution**

**Q.49** Consider the following pseudocode, where S is a semaphore initialized to 5 in line#2 and counter is a shared variable initialized to 0 in line#1. Assume that the increment operation in line#7 is not atomic.

```
1. int counter = 0;
2. Semaphore S = init(5);
3. void parop(void)
4. {
5.     wait(S);
6.     wait(S);
7.     counter++;
8.     signal(S);
9.     signal(S);
10. }
```

If five threads execute the function parop concurrently, which of the following program behavior(s) is/are possible?

- (a) The value of counter is 1 after all the threads successfully complete the execution of parop.
- (b) The value of counter is 0 after all the threads successfully complete the execution of parop.
- (c) There is a deadlock involving all the threads.
- (d) The value of counter is 5 after all the threads successfully complete the execution of parop.

**Ans. (a, c, d)**

(a) True: ( $P_1$ ) first process came executed wait(S); wait(S) counter++; was partially executed and the process preempted. Other process came updated counter to 4.

Now  $P_1$  came and wrote counter to 1. So (a) is correct.

(b) Not possible at least counter will be 1.

(c) True: All 5 processes executed first wait(S) and are blocked.

(d) True: Each process executed the code sequentially and counter was finally updated to 5.

Option (a), (c) and (d) are correct.

**End of Solution**

**Q.50** Which of the following standard C library functions will always invoke a system call when executed from a single-threaded process in a UNIX/Linux operating system?

- (a) exit
- (b) strlen
- (c) sleep
- (d) malloc

**Ans. (a, c)**

**End of Solution**

**Q.51** Consider the following three functions:

$$f_1 = 10^n, f_2 = n^{\log n}, f_3 = n^{\sqrt{n}}$$

Which one of the following options arranges the functions in the increasing order of asymptotic growth rate?

- (a)  $f_2, f_3, f_1$  (b)  $f_2, f_1, f_3$   
(c)  $f_1, f_2, f_3$  (d)  $f_3, f_2, f_1$

**Ans. (a)**

$$\begin{array}{ll} f_2 : n^{\log n} & f_3 : n^{\sqrt{n}} \\ \log(n^{\log n}) & \log(n^{\sqrt{n}}) \\ = (\log n) \cdot (\log n) & = \sqrt{n} \cdot \log n \end{array}$$

$$\therefore n^{\log n} = o(n^{\sqrt{n}}) \quad (f_2 < f_3)$$
$$f_1 = 10^n \leftarrow \text{exponential function}$$

$$\therefore f_1 > f_3 > f_2$$

In increasing order of asymptotic growth rate.

**End of Solution**

**Q.52** Consider the following statements:

$S_1$  : The sequence of procedure calls corresponds to a preorder traversal of the activation tree.

$S_2$  : The sequence of procedure returns corresponds to a postorder traversal of the activation tree.

Which one of the following options is correct?

- (a)  $S_1$  is false and  $S_2$  is false (b)  $S_2$  is true and  $S_2$  is false  
(c)  $S_1$  is true and  $S_2$  is true (d)  $S_1$  is false and  $S_2$  is true

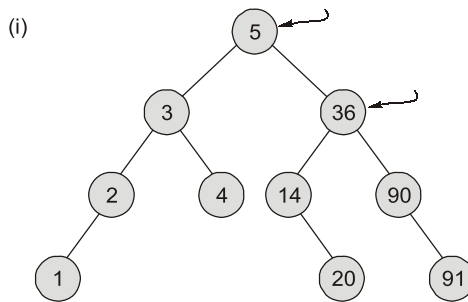
**Ans. (c)**

**End of Solution**

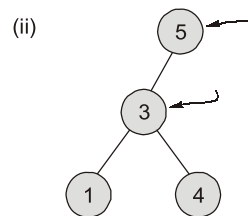
**Q.53** A binary search tree  $T$  contains  $n$  distinct elements. What is the time complexity of picking an element in  $T$  that is smaller than the maximum element in  $T$ ?

- (a)  $\theta(n)$  (b)  $\theta(n \log n)$   
 (c)  $\theta(\log n)$  (d)  $\theta(1)$

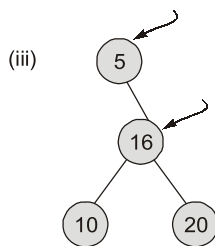
**Ans. (d)**



$\Rightarrow$  5 is definitely lesser than max element



$\Rightarrow$  3 is definitely lesser than max element



$\Rightarrow$  5 is definitely lesser than max element



No such element

In any case, we can find one element that is less in just 2 seeks  $\Rightarrow \theta(1)$ .

**End of Solution**

**Q.54** Consider the following Boolean expression:

$$F = (X + Y + Z)(\bar{X} + Y)(\bar{Y} + Z)$$

Which of the following Boolean expressions is/are equivalent to  $F$ ?

- (a)  $X\bar{Y} + Y\bar{Z} + \bar{X}Y\bar{Z}$  (b)  $(\bar{X} + \bar{Y} + \bar{Z})(X + \bar{Y})(Y + \bar{Z})$   
 (c)  $(X + \bar{Z})(\bar{Y} + \bar{Z})$  (d)  $X\bar{Y} + \bar{Z}$

Ans. (a, c, d)

$$F = (X+Y+Z)(\bar{X}+Y)(\bar{Y}+Z)$$

$$\bar{F} = \overline{(X+Y+Z)(\bar{X}+Y)(\bar{Y}+Z)}$$

$$= \overline{(X+Y+Z)} + \overline{(\bar{X}+Y)} + \overline{(\bar{Y}+Z)} \quad \text{Using Demorgan's theorem}$$

$$= \bar{X}\bar{Y}\bar{Z} + X\bar{Y} + Y\bar{Z}$$

Option (a) is matching.

$$\bar{F} = \bar{Y}(\bar{X}\bar{Z} + X) + Y\bar{Z}$$

$$= \bar{Y}[(\cancel{\bar{Y}} + \bar{X})] (X + \bar{Z}) + Y\bar{Z}$$

$$= \bar{Y}[X + \bar{Z}] + Y\bar{Z}$$

$$= X\bar{Y} + \bar{Y}\bar{Z} + Y\bar{Z}$$

$$= X\bar{Y} + \bar{Z}(\cancel{\bar{Y}} + Y)$$

$$= X\bar{Y} + \bar{Z}$$

Option (d) is matching.

$$X\bar{Y} + \bar{Z} = (X + \bar{Z})(\bar{Y} + \bar{Z}) \quad \text{using distributing property.}$$

So option (c) is matching.

∴ So answer is (a), (c), (d).

End of Solution

**Q.55** Consider the following representation of a number in IEEE 754 single-precision floating point format with a bias of 127.

$$S : 1 \quad E : 10000001 \quad F : 111100000000000000000000$$

Here  $S$ ,  $E$  and  $F$  denote the sign, exponent and fraction components of the floating point representation.

The decimal value corresponding to the above representation (rounded to 2 decimal places) is \_\_\_\_\_.

Ans. (-7.75)

$$\begin{aligned} \text{Value:} & \quad (-1)^S (1.M) \times 2^{E-\text{Bias}} \\ & \quad (-1)^1 (1.11100\dots) \times 2^{10000001 - 127} \\ & \quad -(1.1111) \times 2^{129 - 127} \\ & \quad -1.1111 \times 2^2 \\ & \quad -111.11 \\ & \quad (-7.75)_{10} \end{aligned}$$

End of Solution

